

RSA ENCRYPTION AND PRIME FACTORIZATION

People have needed to communicate information secretly for almost as long as we've been around. We can easily see how this can benefit finance or military, but it's even used in our day-to-day as computers communicate with each other. The earliest form of cryptography used what are known as **symmetric-key ciphers**, where two parties had access to a secret key that could both encrypt and decrypt messages. Of course, this requires the parties to have a way to communicate secretly in the first place. As technology advanced, the need for more sophisticated methods became necessary.

The RSA Cryptosystem—named after Ron Rivest, Adi Shamir, and Len Adleman, the first to publish¹ this method—is what is known as a **asymmetric-key cipher**, where everyone is allowed to encrypt with the public key, but only the holder of the private key can decrypt, making it great for one-way communications! While relatively new, it is built on notions, theorems, and work that has long existed in mathematics (we've covered most of it in class!).

RECALL: The **unit group** of n is the set $\mathbb{Z}_n^\times := \{a \in \mathbb{Z}_n \mid a \text{ is a unit in } \mathbb{Z}_n\}$.

RECALL: Euler's phi function satisfies the following properties:

- (1) If p is prime and n is a positive integer, then $\phi(p^n) = p^{(n-1)}(p - 1)$.
- (2) If m, n are positive coprime integers, then $\phi(mn) = \phi(m)\phi(n)$.

(1) Generating an RSA Key:

- (a) Let $p = 47$ and $q = 59$. Calculate $n = pq$ and find $\phi(n)$.
- (b) Let $e = 17$. Explain why e has an inverse modulo $\phi(n)$.
- (c) Find $d = e^{-1} \pmod{\phi(n)}$.

(2) Encoding and Encrypting:

- (a) Encode the message "HI" into an integer m by converting the letters into numbers according to the table below and concatenating them in order.²

	A	B	C	D	E	F	G	H
00	01	02	03	04	05	06	07	08
I	J	K	L	M	N	O	P	Q
09	10	11	12	13	14	15	16	17
R	S	T	U	V	W	X	Y	Z
18	19	20	21	22	23	24	25	26

- (b) Find $y \equiv m^e \pmod{n}$.

(3) Decoding and Decrypting:

- (a) Find $x \equiv y^d \pmod{n}$ using any techniques³ from class.
- (b) Decode x into a message by reversing the encoding in (2a).
- (c) Explain why $m^{ed} \equiv m \pmod{n}$.
- (d) Encode the message "CAT" as an integer m , then find and compare $y \equiv m^{17} \pmod{2773}$ and $x \equiv y^{157} \pmod{2773}$. Explain why $x \neq m$.

(4) Creating your own key-pair:

- (a) Choose two large primes and compute $n = p \cdot q$ and $\phi(n)$.
- (b) Choose any $0 < e < \phi(n)$ in \mathbb{Z}_n^\times .
- (c) Write your n and e on the board; these make up your public key.
- (d) Find $d = e^{-1} \pmod{\phi(n)}$.

¹Clifford Cocks, an English mathematician, had actually developed a version of this four years prior, but he didn't think it was worth publishing!

²For example, "DOG" becomes $041507 = 41507$.

³HINT: Try using the Chinese Remainder Theorem to work with smaller numbers.

(5) Sending messages⁴:

- (a) Find another group to exchange messages with. Come up with a message m and encrypt it using that group's n and e . Write your encrypted message on the board.
- (b) Once the other group has written their encrypted message for you on the board, decrypt it and see what they sent.
- (c) Pick any group's message on the board and see if you can decrypt it, using any techniques. What do you need to know before you can decrypt the message?

FACTORING METHODS

(6) Factoring by **Trial Division**:

- (a) Let $n = 1643$ be the product of two primes. Factor n by brute force, i.e., attempt to divide by each⁵ prime up to n .
- (b) There is a \$200,000 cash reward for factoring a 617-digit product of two primes. Explain why this is unreasonable to do by Trial Division.

THEOREM: If $a^2 \equiv b^2 \pmod{n}$, then $\gcd(a+b, n) \cdot \gcd(a-b, n) = n$. Furthermore, if $a \not\equiv \pm b \pmod{n}$, then $\gcd(a+b, n)$ and $\gcd(a-b, n)$ are non-trivial factors of n .

(7) Factoring by the **Continued Fraction Algorithm**:

- (a) Let $n = 3053$ be the product of two primes. Find the **factor base** of n : the set of positive primes⁶ $q_i \leq 7$ where $\left(\frac{n}{q_i}\right) = 1$.
- (b) Check⁷ that each element in the factor base is not a prime factor of n .
- (c) Find the first⁸ 5 convergents $C_k = \frac{p_k}{q_k}$ of \sqrt{n} . For each of these, compute $a_k \equiv p_k \pmod{n}$ and $b_k \equiv p_k^2 \pmod{n}$.
- (d) Write each b_k as a product of primes in the factor base, if possible⁹. Find a nonempty set of pairs $(a_i, b_i), \dots, (a_j, b_j)$ such that $b_i \cdots b_j$ is trivially a square modulo n and

$$a_i \cdots a_j \not\equiv \pm \sqrt{b_i \cdots b_j} \pmod{n}$$

- (e) Let $A \equiv a_i \cdots a_j \pmod{n}$ and $B \equiv \sqrt{b_i \cdots b_j} \pmod{n}$. Calculate and compare $A^2 \pmod{n}$ and $B^2 \pmod{n}$.
- (f) Apply the Theorem, and use the Euclidean Algorithm to find the prime factors of n .

⁴If at any point you're waiting, work ahead on future problems!

⁵HINT: Start by determining a reasonable upper bound for the smallest prime factor of n , and then divide and conquer.

⁶The upper bound of 7 was not arbitrary; $7 = \lfloor e^{\frac{1}{2}\sqrt{\ln(n)\ln(\ln(n))}} \rfloor$.

⁷If an element were to be a factor of n , then we can reduce n by that factor and try again.

⁸This choice was arbitrary. If we wish to do this in general, we'll take one convergent at a time until we find a solution.

⁹If b_k isn't possible, try $-b_k = (-1)p_1^{e_1}p_2^{e_2}\cdots p_k^{e_k}$.

(SOME) PRIMES BETWEEN 1000 AND 9999

8539	5801	3251	7487	3083	8269	5749	4127	3823	1871	1567	1777	5711
7817	3529	2239	2797	6691	6247	2579	1307	2749	5813	6091	5651	1499
4337	7589	6143	8741	9283	1321	8011	2657	7043	8369	7219	2311	7681
8629	8039	1097	5021	7561	1237	2161	8849	9467	6571	2741	4549	4421
9533	9391	7793	9007	5849	9479	3643	6053	8171	9209	2069	2459	7193
7159	5501	7841	7573	9859	2647	6679	3163	7649	6173	5011	7541	1291
6277	1789	2609	8573	5303	5657	5179	9649	6131	1753	4597	5953	9551
8761	3881	3407	2699	4493	8677	4657	1481	4457	9629	8599	8147	7549
4591	9067	5479	4229	8819	6287	3637	5927	7247	1217	6421	8093	6427
1439	5557	7451	6529	6491	8287	5323	7753	9103	1033	1609	3187	3023
6911	1973	7457	7499	1913	8123	1901	5851	7879	2693	1259	1123	8467
8689	9161	6737	9239	2969	4729	6899	2131	7919	9419	3391	9511	2113
3931	2753	1487	9967	1367	7703	3079	1069	6121	4157	1549	6043	9371
9721	8821	9199	9491	1627	3607	4139	1427	1399	9187	7949	9397	6373
9203	2383	5861	6547	4651	8543	5189	5683	6833	1999	1453	5413	6823
5449	3709	6971	5669	6977	7591	1301	8753	6197	2269	8389	6367	7547
1153	8867	3793	2557	6151	8537	3181	9803	5807	3049	4751	9949	9403
8209	1733	1741	3923	7019	3119	4639	8647	5113	3697	2879	7907	5923
5233	3889	9781	4523	5023	7993	8087	9539	4793	2437	9839	8641	3329
1669	2251	4801	3659	4049	9767	1093	2617	5417	8191	4261	3581	3673
5903	4783	6983	7013	8713	9349	6079	6599	3037	7621	6857	9241	2777
7243	9973	7207	6959	1229	7351	5641	6569	7433	5387	9311	5527	9689
3121	2399	3547	1091	6553	1279	1721	2273	7607	8951	7529	6073	7727
7901	1759	2423	3539	4831	2333	3089	7321	7669	6163	4759	4691	4969
8513	9857	3911	1289	1201	3853	5843	5647	7109	3677	3217	4519	1277
3767	3833	7757	5689	1979	1483	8837	3313	9787	9337	3449	4663	5653
1879	8563	9437	2153	5779	2347	8923	5279	4211	8419	6067	5419	9227
1621	9011	9619	6709	4703	5167	8971	9041	4463	3671	5347	8803	8429
3727	8069	1543	3533	3739	7507	2591	1559	6761	8831	1447	4093	6961
3469	9473	7489	1787	3343	4363	7829	8317	4679	4051	6359	7349	4021
7309	8297	5879	2099	9059	6011	7537	6449	4243	9431	4909	9341	6829
7213	2417	1187	8719	6997	4153	7297	9377	9293	9631	6089	2441	4253
1423	1451	4111	9091	1723	5273	6217	1009	2917	2897	1303	9661	8293
2677	6133	1061	1213	7883	5003	7723	7517	7691	1571	5443	5581	7307
3613	8969	4973	8929	3323	1063	2137	2833	4273	3433	2857	4129	1373
5399	4957	2539	2287	3701	5981	2887	7027	4583	4231	1747	4567	2039
1019	1021	2521	3557	4007	9871	3299	1049	5039	9497	1163	2473	2927
4507	7477	9221	8053	8807	4219	8081	3221	7417	2711	3373	8623	3851
7717	3821	7577	2063	1231	4861	7393	8009	7699	5087	3389	6113	9181
7069	6221	6451	9109	7559	6703	4817	9851	2551	2341	4721	2861	3907
2203	4391	3011	8747	3623	7283	6871	3359	8311	9749	7411	2351	4951
3361	3919	2027	8219	6361	9613	4919	2281	5821	8443	6803	1429	6299
8839	4517	6661	2801	8447	5147	6397	5531	1867	5119	1583	3041	9739
1601	4649	6389	4073	7867	1693	6883	3541	8779	9043	8017	3467	2339
8461	5297	9343	2357	1531	6791	9929	3691	6733	8933	2971	9817	8273
4079	1129	8699	6203	1523	4673	6907	5869	2803	3229	2707	5659	7369
4483	8179	8893	6779	6563	8731	2659	3929	3559	3527	2687	3779	4877
9461	8861	4177	4003	1283	4733	1223	8291	7103	1667	7853	5227	4993